

# On consistent families of criteria : An axiomatic approach

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# Outline

- Introduction: Consistent families in the sense of Roy
- Example: Ranking vectors of constraint satisfaction degrees
- A general model
  - Characterisation
  - Interpretation
  - Back to the example
- Discussion and future work

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# Introduction

## Families of criteria (Roy, 1985)

- Definition of a criterion:  $g_i : \mathcal{A} \rightarrow \mathbb{R}$   
interpretation of  $g_i(a) \geq g_i(b)$ :  $a$  is at least as good as  $b$   
on point of view  $i$
- $n$  points of view:  $i = 1, \dots, n$
- Notation for the global preference on  $\mathcal{A}$  (to be constructed by applying some method):  $\succsim$

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# Consistent families

## Properties to be fulfilled

- Exhaustivity:

$$[g_i(a) = g_i(b), \forall i] \Rightarrow a \sim b$$

- Consistency:

$$\left. \begin{array}{l} a \succsim b \\ g_i(c) \geq g_i(a) \quad \forall i \\ g_i(b) \geq g_i(d) \quad \forall i \end{array} \right\} \Rightarrow c \succsim d$$

- Non-redundancy

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# Interpretation and Question

## Interpretation ...

... of consistency : The relationship between the criteria and the constructed preference  $\succsim$  is supposed to be monotonic

**Alternative formulation :** respect of dominance

## Question :

Is it possible to characterise the preferences that are consistent with a family of criteria ?

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# Example

## Comparing vectors of satisfaction levels

### **FCSP** : Flexible Constraint Satisfaction Problem

- vertices = tasks
- arcs = constraints  $C_i, i = 1, \dots, n$
- solution = a feasible schedule for the tasks
- quality of solution : a vector of satisfaction levels

$$x = (x_1, \dots, x_n)$$

- problem : comparing the solutions

General problem of “vector optimisation” : no evident complete ordering of the solutions

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## Dominance

A natural partial order : the dominance relation  $D$  (Pareto ordering)

$$x D y \text{ if } x_i \geq y_i, \forall i$$

Note: in general  $x_i$  belongs to a set  $X_i$  that is at least ordered

**Question :** Characterise “natural and operational methods” leading to a complete ordering of the vectors and preserving dominance (question raised by Dubois and Prade)

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## A framework

- $X = X_1 \times X_2 \times \dots \times X_n$  : finite set of alternatives
- $N = \{1, 2, \dots, n\}$ : set of attributes
- $x = (x_1, x_2, \dots, x_n) \in X$
- $x_{-i} = (x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n)$
- $(a_i, x_{-i}) \in X$
- $\succsim$  reflexive binary relation on  $X$  interpreted as “large preference”
- $\sim$  : symmetric part;  $\succ$  : asymmetric part

Note : there is no structure assumed on  $X_i$



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## Model 0

$$x \succsim y \Leftrightarrow F((u_i(x_i), u_i(y_i)), i = 1, \dots, n) \geq 0$$

for some functions  $u_i : X_i \rightarrow \mathbb{R}$

and some function  $F : \mathbb{R}^{2n} \rightarrow \mathbb{R}$

Note : no properties assumed on functions  $F$  and  $u_i$

**Result :** any relation  $\succsim$  on  $X$  satisfies Model 0

Note :  $u_i$  induces an ordering  $\succsim_i^\pm$  on  $X_i$  :

$$x_i \succsim_i^\pm y_i \Leftrightarrow u_i(x_i) \geq u_i(y_i)$$

This ordering plays no role (is arbitrary) when  $F$  has no special property

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## Model 1

$$x \succsim y \Leftrightarrow F((u_i(x_i), u_i(y_i)), i = 1, \dots, n) \geq 0$$

for some functions  $u_i : X_i \rightarrow \mathbb{R}$   
and some function  $F : \mathbb{R}^{2n} \rightarrow \mathbb{R}$   
such that

$$F(\nearrow, \searrow)$$

i.e.  $F$  is non-decreasing in its first  $n$  coordinates  
and non-increasing in its last  $n$  coordinates:

**Property :**  $\succsim$  is monotonic (consistent) with respect to the orders  $\succsim_i^\pm$

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$$\left. \begin{array}{l} x \succsim y \\ u_i(z_i) \geq u_i(x_i) \quad \forall i \\ u_i(y_i) \geq u_i(w_i) \quad \forall i \end{array} \right\} \Rightarrow z \succsim w$$

Indeed :

$$\left. \begin{array}{l} F(u_i(x_i), u_i(y_i)) \geq 0 \\ F(\nearrow, \searrow) \end{array} \right\} \Rightarrow F(u_i(z_i), u_i(w_i)) \geq 0$$

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## Characterisation of Model 1

**Result :** the relations that satisfy Model 1 are those verifying properties AC123

**AC1<sub>i</sub> :**

$$\left. \begin{array}{c} (x_i, a_{-i}) \succsim (y_i, b_{-i}) \\ \text{and} \\ (z_i, c_{-i}) \succsim (w_i, d_{-i}) \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} (z_i, a_{-i}) \succsim (y_i, b_{-i}) \\ \text{or} \\ (x_i, c_{-i}) \succsim (w_i, d_{-i}), \end{array} \right.$$

**Definition of  $\succsim_i^+$  :**

$$x_i \succsim_i^+ y_i \Leftrightarrow [(y_i, a_{-i}) \succsim z \Rightarrow (x_i, a_{-i}) \succsim z]$$

**Result 1 :** AC1<sub>i</sub>    iff     $\succsim_i^+$  is a complete preorder

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**Result 2 :**  $AC2_i$  iff  $\succsim_i^-$  is a complete preorder  
with

$$x_i \succsim_i^- y_i \Leftrightarrow [z \succsim (x_i, a_{-i}) \Rightarrow z \succsim (y_i, a_{-i})],$$

**Result 3 :**  $AC123_i$  iff  $\succsim_i^\pm$  is a complete preorder  
with

$$x_i \succsim_i^\pm y_i \Leftrightarrow [x_i \succsim_i^+ y_i \text{ and } x_i \succsim_i^- y_i],$$

### Consequence :

If  $\succsim$  satisfies  $AC123$ ,

- the relations  $\succsim_i^\pm$  are complete preorders
- $u_i$  can be chosen to be any numerical representation of  $\succsim_i^\pm$
- $F(u_i(x_i), u_i(y_i))$  is easily built

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## Model 2

A model for a preference  $\succsim$  that is a **complete preorder** :

$$\begin{aligned} x \succsim y &\Leftrightarrow u(x) = u(y) \\ &\Leftrightarrow G(u_i(x_i), i = 1, \dots, n) \geq G(u_i(y_i), i = 1, \dots, n) \end{aligned}$$

for some functions  $u_i : X_i \rightarrow \mathbb{R}$

and some function  $G : \mathbb{R}^n \rightarrow \mathbb{R}$

Model 2  $\Leftrightarrow$  Model 0 with :

$$F(u_i(x_i), u_i(y_i)) = G(u_i(x_i), i = 1, \dots, n) - G(u_i(y_i), i = 1, \dots, n)$$

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## Model $1 \wedge 2$

**Result :**

$\succsim$  is a complete preorder and satisfies AC123

$\Leftrightarrow$

$$\exists u_i : X_i \rightarrow \mathbb{R}, \quad \exists G : \mathbb{R}^n \rightarrow \mathbb{R}$$
$$G(\nearrow) \text{ (non-decreasing)}$$

such that

$$x \succsim y \Leftrightarrow G(u_i(x_i), i = 1, \dots, n) \geq G(u_i(y_i), i = 1, \dots, n)$$

**Examples :**

- $G = \sum_i u_i(x_i)$  (MAU)
- $G = \max_i \{u_i(x_i)\}$

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## Back to the example FCSP

... and coherent families

What if the  $X_i$  are ordered *a priori* ?

If a preference  $\succsim$  is built that is a complete preorder and respects dominance, then the relations  $\succsim_i^\pm$  that can be deduced from  $\succsim$  are compatible with the complete preorders given *a priori*.



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## Conclusions and future research

- Framework in which specific methods can be further specified (max, MAU, ... )
- There are further interesting “sub frameworks” : e.g. the case where  $F$  (or  $G$ ) is **strictly** monotonic (already characterised)
- Advantages of such a framework :
  - understanding
  - suggests methods of elicitation
  - allows for choosing between several specific procedures in a “family”