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NONCOMPENSATORY AND GENERALIZED NONCOMPENSATORY PREFERENCE STRUCTURES

1. INTRODUCTION

The aim of this paper is to provide a general study of noncompensatory preference structures. These structures have not been studied very much yet, the attention of most decision theorists being almost exclusively devoted to structures allowing some kind of utility representation. They nevertheless appear frequently in practice both as heuristic approaches to analyse multidimensional evaluations (e.g. disjunctive and lexicographic models, see MacCrimmon, 1973) and as easy to implement methods to perform an aggregation of several attributes for decision-aid (ELECTRE methods, see Roy, 1968, 1971; Roy and Bertier, 1972).

The paper is organized as follows. We present our notations in Section 2. In Section 3 we recall some definitions and propositions about noncompensatory preference structures and introduce the notion of concordance preference structure. In Section 4 we propose a generalization of these notions introducing the idea of discordance. Section 5 provides a brief description of how such preference structures can be used for decision-aid.

2. NOTATIONS AND PRELIMINARY DEFINITIONS

Throughout the paper we will note:
\( \mathbb{R}_+^* \) the set of strictly positive real numbers,
\( \mathbb{N} = \{0, 1, 2, \ldots\}, \mathbb{N}_0 = \mathbb{N} \setminus \{0\}, \)
\( \Omega = \{1, 2, \ldots, n\} \) with \( n \in \mathbb{N} \) and \( n \geq 2, \)
\( P(\Omega) \) the set of all subsets of \( \Omega, \)
\( S \) the set of all pairs of disjoint subsets of \( \Omega, \)
\( S = \{(A, B) \mid A, B \in P(\Omega) \text{ and } A \cap B = \emptyset\}, \)
\( X_1, X_2, \ldots, X_n \) nonempty sets which can be interpreted as \( n \) sets of levels defining \( n \) attributes in a multi-attribute decision problem,

*Theory and Decision* 21 (1986) 251–266.
\( X = X_1 \times \ldots \times X_n \), the Cartesian product of these sets.

\((x_i, (y_j)_{j \neq i})\) the element of \( X \) \((y_1, y_2, \ldots, y_{i-1}, x_i, y_{i+1}, \ldots, y_n)\).

\(\succ\) an asymmetric binary relation on \( X \) which can be interpreted as a (strict) preference relation.

\((X, \succ)\) will be called a Preference Structure (P.S.).

We will classically note \(\sim\) the binary relation on \( X \) defined by \(x \sim y\) iff not \((x \succ y)\) and not \((y \succ x)\) and \(\succeq\) the binary relation on \( X \) such that \(x \succeq y\) iff not \((y \succ x)\).

**DEFINITION 2.1.** For all \(i \in \Omega\) we define a binary relation \(\succ_i\) on \(X_i\) by

\[ x_i \succ_i y_i \text{ iff } (x_i, (a_j)_{j \neq i}) \succ (y_i, (a_j)_{j \neq i}) \text{ for all } (a_j)_{j \neq i} \in X_{j \neq i} \times X_{j i}. \]

From \(\succ_i\) we define \(\sim_i\) and \(\succeq_i\) as above.

The asymmetry of \(\succ\) obviously implies the asymmetry of each \(\succ_i\). The definition of \(\succ_i\) does not imply any notion of preferential independence since we have \(x \succ y\) only if \((x, (a_j)_{j \neq i}) \succeq (y, (a_j)_{j \neq i})\) for all \(j \neq i\).

**DEFINITION 2.2.** An attribute \(X_i\) is essential iff \(x_i \succ_i y_i\) for some \(x_i, y_i \in X_i\).

We will assume hereafter that all attributes are essential which will prove unrestrictive for our purposes.

**DEFINITION 2.3.** For each ordered pair \((x, y) \in X^2\) we will note

\[ P(x, y) = \{i \in \Omega \mid x_i \succ_i y_i\} \]

Thus \(P(x, y)\) denotes the set of attributes for which there is a partial preference for \(x\) on \(y\). The asymmetry of each \(\succ_i\) implies that

\[ P(x, y) \cap P(y, x) = \emptyset \text{ for all } x, y \in X. \]

**DEFINITION 2.4.** We will note \(\triangleright\) and \(=\) the binary relations on \(P(\Omega)\) defined respectively by \(A \triangleright B\) iff \((P(x, y), P(y, x)) = (A, B)\) for some \(x, y \in X\).

It is clear that \(A \triangleright B\) implies \((A, B) \in S\) and that the following lemma holds for all \(A, B \in P(\Omega)\).

**LEMMA. 2.1.** \(A = B \iff A \cap B = \emptyset\) and \(\forall x, y \in X: (P(x, y), P(y, x)) = (A, B) = x - y\).
DEFINITION 2.5. We will note $\succ$ the binary relation on $P(\Omega)$ defined by: $A \succ B$ iff $A \cap B = \emptyset$ and $\forall x, y \in X$: $(P(x, y), P(y, x)) = (A, B) \Rightarrow x \succ y$.

Contrary to $\triangleright$, $\succ$ is asymmetric when all attributes are essential. According to Definition 2.5, $\succ$ can be interpreted as a "more important than" relation on $P(\Omega)$.

- $P_1$ superadditivity iff $((A \cup C) \cap (B \cup D)) = \emptyset$ and $A \triangleright B$ and $C \triangleright D$, $A \cup C \triangleright B \cup D$.
- $P_2$ decisivity iff $((A, B) \in S$ and $(A, B) \neq (\emptyset, \emptyset)) \Rightarrow (A \sim B)$.
- $P_3$ attribute acyclicity iff $\triangleright$ has no cycles.
- $P_4$ attribute transitivity iff $A \triangleright B, B \triangleright C$ and $A \cap C = \emptyset \Rightarrow A \triangleright C$.
- $P_5$ double essentiality iff $\forall i \in \Omega, x_i \succ y_i$ and $y_i \succ z_i$ for some $x_i, y_i, z_i \in X_i$.

3. NONCOMPENSATORY PREFERENCE STRUCTURES (NPS)

In this section we first recall some definitions and propositions about the notion of noncompensatory preference structures introduced independently by Fishburn (1974, 1975 and 1976) and Plott et al. (1975). Two special important cases of NPS are analysed: the classical lexicographic preferences and the new concept of concordance preference structures.

DEFINITION 3.1. A P.S. $(X, \succ)$ is noncompensatory iff $\forall x, y, z, w \in X$:

(1) $[(P(x, y), P(y, x)) = (P(z, w), P(w, z))] \Rightarrow (x \succ y \Rightarrow z \succ w)$.

(2) $P(x, y) \neq \emptyset$ and $P(v, y \setminus v) \neq \emptyset \Rightarrow y \succ v$.

The idea of noncompensation appears clearly in this definition since the global preference of $x$ on $y$ only depends on the subsets of $\Omega$ on which there is a partial preference of $x$ on $y$ and of $y$ on $x$. This definition corresponds to a "regular noncompensatory preference structure" in Fishburn (1976). Condition (2) of Definition 3.1. could be omitted, but from a practical point of view only regular structures are of interest.

It results immediately from Definition 3.1. that if $(X, \succ)$ is a NPS then the attributes are mutually preferentially independent (cf. Keeney and Raiffa, 1976) (which can be shown to be implied by condition (1) alone). We have:
LEMMA 3.1. A P.S. \((X, \succ)\) is a NPS iff:
(1) \(\forall A, B \in P(\Omega), A \triangleright B \Rightarrow A \triangleright B\).
(2) \(\forall A \in P(\Omega) \setminus \emptyset, A \triangleright \emptyset\).

Proof. Left to the reader.

Condition (1) of Lemma 3.1. obviously implies that for a NPS\(\triangleright\) is asymmetric and that \(\emptyset \approx \emptyset\).

Lexicographic preference structures are an important particular case of NPS.

DEFINITION 3.2. A P.S. \((X, \succ)\) is lexicographic iff there is a permutation \(\sigma\) on \(\Omega\) such that \(\forall x, y \in X, x \succ y\) iff not \((x_i \sim_i y_i)\) for some \(i \in \Omega\) and \(x_{\sigma(i)} \succ_{\sigma(i)} y_{\sigma(i)}\) for the smallest \(i\) for which not \((X_{\sigma(i)} \sim_{\sigma(i)} Y_{\sigma(i)})\).

Fishburn (1976) Theorem 1, proves the following results:

THEOREM 3.1. If \((X, \succ)\) is a NPS, then:
(a) \(P_2\) and \(P_3\Rightarrow P_4\),
(b) \(P_1\) and \(P_2\) and \(P_3\Rightarrow (X, \succ)\) is lexicographic,
(c) \(P_3\) and \(\succ\) is a weak order \(\Rightarrow (X, \succ)\) is lexicographic.

The notion of NPS also provides a new insight into the idea of concordance which appears in a wide variety of multicriteria decision-aid methods such as ELECTRE. We formalize here this notion using an idea introduced by Vansnick (1986). For another approach to the idea of concordance we refer to Huber (1974 and 1979).

DEFINITION 3.3. A P.S. \((X, \succ)\) is a concordance P.S. of type \(\rho\) (CPS\(\rho\)) with \(\rho \geq 1\) a rational number, iff
\[
\exists f_1, f_2, \ldots, f_n \in P(\Omega) \text{ such that } \forall x, y \in X, x \succ y \Leftrightarrow \sum_{i \in P(x, y)} f_i > \rho \sum_{j \in P(y, x)} f_j.
\]

Figure 1 gives a graphical interpretation of this definition.
We have the following lemmas:

LEMMA 3.2. \(\forall \rho \geq 1, \text{ a CPS}\(\rho\) is a NPS verifying \(P_4\).

Proof. Obvious (left to the reader).
LEMMMA 3.3. A lexicographic NPS is a CPS1.

Proof: Follows immediately from taking $f_{0(0)} = 2^{(n+1) - 1}$.

The following theorem gives necessary and sufficient conditions for the existence of a CPSp. These conditions are very similar to those appearing in generalizations of compensatory models (Domotor and Stelzer, 1971, Fishburn, 1969; Krantz et al., 1971, chap. 9; Roberts, 1979, chap. 8). This is not surprising because the relation “more important than” between attributes has strong connections with a “more probable than” relation between events. In this theorem, $\forall j, y \in X$, $M(x, y)$ will denote the $1 \times n$ matrix $(a_1, a_2, \ldots, a_n)$ where $\forall i \in \Omega$, $a_i = 1$ iff $x_i >_i y$ and $a_i = 0$ otherwise.

THEOREM 3.2. A P.S. $(X, \succ)$ is a CPSp iff:

$(X, \succ)$ is a NPS and

$$\forall m, k \in \mathbb{N}_0$$

$$\sum_{i=1}^{\infty} M(x^{(i)}, y^{(i)}) + \rho \sum_{0 \leq j \leq \infty} M(z^{(j)}, w^{(j)}) \neq$$
\[ (3.2) \quad \rho \sum_{i=1}^{m} M(y^{(i)}, x^{(i)}) + \sum_{0<j<k} M(w^{(j)}, z^{(j)}) \]

whenever

\[ x^{(i)}, y^{(i)} \in X \text{ and } x^{(i)} \succ y^{(i)} \quad \forall \quad i \in [1, 2, \ldots, m] \]

\[ z^{(j)}, w^{(j)} \in X \text{ and } w^{(j)} \sim z^{(j)} \quad \forall \quad j \in \mathbb{N} \text{ such that } 0<j<k. \]

Let us notice that condition (3.2) of Theorem 3.2. contains in fact an infinity of conditions which are rather delicate to interpret as this is the case for the theory of comparative probability. We refer to Vansnick (1986) for a more general presentation of a similar theorem including a threshold.

**Proof.** Theorem 3.2. can be stated:

\[ \exists F = \begin{pmatrix} f_1 \\ \vdots \\ f_n \end{pmatrix} \in [\mathbb{R}_0^+]^{n+1} \text{ such that, } \forall \ x, y \in X \]

\[ x \succ y \Rightarrow M(x, y) \cdot F \geq \rho \ M(y, x) \cdot F \]

\[ x \sim y \Rightarrow M(x, y) \cdot F \leq \rho \ M(y, x) \cdot F \]

iff (3.1) and (3.2).

(a) **Necessity**

The necessity of (3.1) is obvious. Let \( m, k \in \mathbb{N}_0 \), \( x^{(i)} \succ y^{(i)} \quad \forall \quad i \in [1, 2, \ldots, m] \) and \( w^{(j)} \sim z^{(j)} \quad \forall \quad j \in \mathbb{N} \text{ such that } 0<j<k. \)

By assumption, there is a \( F \in [\mathbb{R}_0^+]^{n+1} \) such that

\[ M(x^{(i)}, y^{(i)}) \cdot F \geq \rho \ M(y^{(i)}, x^{(i)}) \cdot F \quad \forall \quad i \in [1, 2, \ldots, m] \]

\[ \rho \ M(z^{(j)}, w^{(j)}) \cdot F \geq M(w^{(j)}, z^{(j)}) \cdot F \quad \forall \quad j \in \mathbb{N} \text{ such that } 0<j<k. \]

After summation, we obtain:

\[ [\sum_{i=1}^{m} M(x^{(i)}, y^{(i)}) + \rho \sum_{0<j<k} M(z^{(j)}, w^{(j)})] \cdot F \geq \]

\[ [\rho \sum_{i=1}^{m} M(y^{(i)}, x^{(i)}) + \sum_{0<j<k} M(w^{(j)}, z^{(j)})] \cdot F \]

which implies (3.2).
(b) Sufficiency

Let us observe that it is sufficient to establish that (3.1) and (3.2) imply:

\[ \exists F \in \mathbb{R}^{n+1} \text{ such that, } \forall x, y \in X: \]

\[ x \succ y = M(x, y) \cdot F \succ \rho M(y, x) \cdot F \text{ and } \]

\[ x - y = M(x, y) \cdot F \preceq \rho M(y, x) \cdot F \]

(3.3)

since (3.1) and (3.3) imply that \( f_i > 0 \) \( \forall i \in \Omega \). We will show that (3.1) and not (3.3) = not (3.2). Let \( Y^2 \) be a set containing one element from each equivalence class of the relation \( E \) defined on \( X^2 \) by:

\[ (x, y) \sim (z, w) \text{ iff } P(x, y) = P(z, w) \text{ and } P(y, x) = P(w, z). \]

\( n \) being finite, \( Y^2 \) contains a finite number of elements.

Let \( \{(s^{(i)}, t^{(i)}) \mid i = 1, 2, ..., I\} \) be the set of elements in \( Y^2 \) such that \( s^{(i)} \succ t^{(i)} \).

Let \( \{(u^{(j)}, v^{(j)}) \mid j = 1, 2, ..., J\} \) be the set of elements in \( Y^2 \) such that \( u^{(j)} \sim v^{(j)} \).

Given (3.1), not (3.3) is equivalent to: \( \exists F \in \mathbb{R}^{n+1} \) such that:

\[ M(s^{(i)}, t^{(i)}) \cdot F \succ \rho M(t^{(i)}, s^{(i)}) \cdot F \forall i \in [1, 2, ..., I] \]

and

\[ M(u^{(j)}, v^{(j)}) \cdot F \preceq \rho M(v^{(j)}, u^{(j)}) \cdot F \forall j \in [1, 2, ..., J]. \]

Each attribute being essential, we have \( I \geq n > 0 \). Therefore, according to Motzkin’s transposition theorem (see Vansnick, 1984), there are \( \lambda_1, \lambda_2, ..., \lambda_1, \mu_1, \mu_2, ..., \mu_J \in \mathbb{R}^+ \) with \( \lambda_i \neq 0 \) for at least one \( i \in [1, 2, ..., I] \) such that

\[ \sum_{i=1}^{I} \lambda_i M(s^{(i)}, t^{(i)}) + \rho \sum_{j=1}^{J} \mu_j M(v^{(j)}, u^{(j)}) = \]

\[ \rho \sum_{i=1}^{I} \lambda_i M(t^{(i)}, s^{(i)}) + \sum_{j=1}^{J} \mu_j M(u^{(j)}, v^{(j)}). \]

As the elements of \( M(s^{(i)}, t^{(i)}), M(t^{(i)}, s^{(i)}), M(v^{(j)}, u^{(j)}) \) and \( M(u^{(j)}, v^{(j)}) \) are either 0 or 1 and \( \rho \) is a rational number, there exist \( \lambda_1^*, \lambda_2^*, ..., \lambda_I^*, \mu_1, \mu_2, ..., \mu_J \in \mathbb{N} \) such that
\[ \sum_{i=1}^{i} \lambda_i^* M(v^{(i)}, u^{(i)}) + \rho \sum_{j=1}^{j} \mu_j^* M(v^{(j)}, u^{(j)}) = \rho \sum_{i=1}^{i} \lambda_i^* M(v^{(i)}, s^{(i)}) + \sum_{j=1}^{j} \mu_j^* M(v^{(j)}, s^{(j)}) \]

with

\[ \sum_{i=1}^{i} \lambda_i^* > 0. \]

(3.4)

(1) \[ m = \sum_{i=1}^{i} \lambda_i^* \quad k = \sum_{j=1}^{j} \mu_j^* + 1 \]

(2) for \( i = 1, 2, \ldots, I \) and \( h = 1, 2, \ldots, m \)
\[ x^{(h)} = s^{(i)} \quad \text{and} \quad y^{(h)} = t^{(i)} \text{ if} \]
\[ \sum_{i=1}^{i} \lambda_i^* < h \quad \text{and} \quad h \leq \sum_{i=1}^{i} \lambda_i^* \]

(3) for \( j = 1, 2, \ldots, J \) and \( h' \in \mathbb{N} \) such that \( 0 < h' < k \)
\[ w^{(h')} = u^{(i)} \quad \text{and} \quad z^{(h')} = v^{(i)} \text{ if} \]
\[ \sum_{i=1}^{i} \mu_i^* < h' \quad \text{and} \quad h' \leq \sum_{i=1}^{i} \mu_i^* \]

were by convention \( \mu_0^* = \lambda_0^* = 0 \)

we have:

\[ \forall k, m \in \mathbb{N}_0 \quad x^{(h)} \succ y^{(h)} \quad \forall h \in [1, 2, \ldots, m] \]

\[ w^{(h')} \succ z^{(h')} \quad \forall h' \in \mathbb{N} \] such that \( 0 < h' < k \)

and 3.4. can be written

\[ \sum_{h=1}^{m} M(x^{(h)}, y^{(h)}) + \rho \sum_{0 < h' < k} M(z^{(h')}, w^{(h')}) = \]

\[ \rho \sum_{h=1}^{m} M(y^{(h)}, x^{(h)}) + \sum_{0 < h' < k} M(u^{(h')}, v^{(h')}) \]

which completes the proof.
4. GENERALIZED NONCOMPENSATORY PREFERENCE STRUCTURES (GNPS)

It is often interesting, from a practical point of view, to weaken the absolute noncompensation of NPS in order to obtain more realistic comparisons (see Roy, 1974; Huber, 1979; and Vansnick, 1986). This is the purpose of the following definition:

DEFINITION 4.1. A P.S. \((X, \succ)\) is a GNPS iff \(\forall x, y, z, w \in X:\)
(1) \((P(x, y), P(y, x)) = (P(z, w), P(w, z))\Rightarrow [x \succ y = z \succ w].\)
(2) \([P(x, y) \neq \emptyset \text{ and } P(y, x) = \emptyset] \Rightarrow x \succ y.\)

This definition represents a natural generalization of Definition 3.1. Allowing a value of the same kind \((0, 0), 1\) with \(P(x, 0) \neq 0, P(0, x) \neq 0\).

\(x \succ y\) and \(z \sim w\). The possibility of an absence of preference between \(z\) and \(w\) aims to encompass the notion of discordance between evaluations (see Roy, 1968; and Roy and Bertier, 1972). In fact, when the difference between the evaluations of \(z\) and \(w\) becomes important on the attributes belonging to \(P(w, z)\) it is unrealistic to suppose \(z \succ w\).

Definition 4.1. obviously implies mutual preferential independence in a GNPS. We have:

LEMMA 4.1. A P.S. \((X, \succ)\) is a GNPS iff
(1) \(\succ\) is asymmetric.
(2) \(\forall x \in P(\emptyset) \setminus \emptyset, A \sim \emptyset.\)

Proof. Follows immediately from essentiality and Definition 4.1. 

It is obvious from the definition that a NPS is also a GNPS. The following definition establishes an interesting link between NPS and GNPS.

DEFINITION 4.2. Let \((X, \succ)\) be a GNPS. We define on \(X\) a binary relation \(\succeq\) by \(x \succeq y\) iff \(P(x, y) \succ P(y, x)\).

We have:

LEMMA 4.2. If \((X, \succ)\) is a GNPS, then \(\forall x, y, z \in X:\)
(1) \(x \succ y \Rightarrow x \succ y.\)
(2) \((X, \succeq)\) is a NPS.
Proof. Obvious; left to the reader.

Thus, there is a natural way to extend a GNPS into a NPS. For instance, let \( X = \{x_1, y_1\} \times \{x_2, y_2, z_2\} \). We can represent a P.S. by mean of the following matrix:

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It is easily seen that this P.S. is a GNPS with \( x_1 \succ 1, y_1, x_2 \succ 2, y_2 \), \( x_2 \succ 2, z_2, x_2 \succ 2, z_2 \) and \( 1] \gg \emptyset, [2] \gg \emptyset, [1, 2] \gg \emptyset \) and \( 1] \gg [2] \).

In order to obtain its corresponding NPS, it suffices to add \((x_1, z_2) \succ (v_1, x_2) \) and \((x_1, z_2) \succ (v_1, y_2) \).

Any configuration in the darkened area of the matrix with at least one \( > \) would have lead to the same associated NPS.

The following considerations allow to specify better the way the discordance effect can work in a GNPS.

**Definition 4.3.** A GNPS is discordant iff

\( \forall x, y \in X \) such that \( x \not\succ y \):

\[ \forall j \in P(y, x), \exists (x_i^j), (y_i^j), j \in X \times X \text{ such that } ((x_i^j), x, j) > ((y_i^j), x, j) \Rightarrow x > y. \]

In order to be able to interpret this definition, we will use the following:

**Definition 4.4.** \( \forall j \in \Omega \), we define a binary relation \( V_j \) on \( X_j \) by \( x_j V_j y_j \) iff \( \exists (x_i^j), (y_i^j), j \in X \times X \text{ such that } (y, (y_i^j), x, j) > (x, (x_i^j), x, j). \)
The following lemma establishes the link between these two definitions:

**Lemma 4.3.** In a discordant GNPS

\[ x \succ y \text{ iff } x \succ y \text{ and not } y_j V_j x_j \text{ for all } j \in P(y, x). \]

*Proof.* Follows immediately from Definitions 4.3. and 4.4.

Therefore, in a discordant GNPS, the discordance effect is introduced whenever there is an attribute \( j \) in \( P(y, x) \) for which \( y_j V_j x_j \) which can be interpreted as "\( y_j \) is far better than \( x_j \)". It should be noticed that this definition implies that each attribute must be considered separately in order to decide for the discordance. Thus, there is no possibility of interaction between the attributes in \( P(y, x) \).

We have:

**Lemma 4.4.** \( \forall j \in \Omega \), \( V_j \) is asymmetric

*Proof.* \( x_j \succ y_j \) implies by definition not \( y_j V_j x_j \). Thus \( y_j V_j x_j \) implies:
- either \( x_j \succ y_j \) and \( V_j \) is impossible
- or \( y_j \nless j x_j \). The definition of a GNPS and essentially of attribute \( k \) thus imply \( (v_k, x_j, (v'_j)_k, j) \succ (v_k, x_j, (v'_j)_k, j) \) for some \( (v'_j)_k, j \in X_i \). Thus we cannot have \( x_j V_j y_j \).

In the case where a weak order underlies each \( \succ \) some monotonicity conditions on \( \succ \), give rise to a semi-order structure for each \( V_j \).

Let us first recall (see for instance Vincke, 1980) that if \( \succ \) is a semi-order on \( X_i \), the binary relation \( T_i \) on \( X_i \) defined by:

\[
y_i \succ T_i x_i \text{ iff } \forall z_i \in X_i: \begin{cases} z_i \succ i y_i \Rightarrow z_i \succ i x_i \\ x_i \succ i z_i \Rightarrow y_i \succ i z_i \end{cases}
\]

is transitive and strongly connected.

Its asymmetric part \( T_i^A \) is thus a weak order and its symmetric part \( T_i^S \) an equivalence.

We can now state the following:

**Theorem 4.1.** If a discordant GNPS is such that:

\[
\begin{align*}
\succ j \text{ is a semi-order } \forall j \in \Omega, \\
x \succ y & \iff \forall j \in \Omega, \forall w_j \in X_j \text{ such that } y_j T_j w_j: \\
x \succ (y_j)_i \neq j w_j).
\end{align*}
\]
then each $V_j$ is an interval order. Furthermore, if the GNPS also verifies

\[(4.3) \quad x > y = \{v \in \Omega : \forall j \in X_j \text{ such that } z_j T^4_j x_j : ((x_j)_j \neq y) \} \supset y,\]

then $V_j$ is a semi-order.

In order to prove Theorem 4.1, we need the following:

**Lemma 4.5.** Under the assumptions (4.1) and (4.2), we have $\forall j \in \Omega$, $y_j \not V_j x_j$ iff $y_j T^4_j z_j$ for all $z_j \in S(x_j)$ where $S(x_j) = \{y_j \in X_j : (x_j)_j \neq x_j \}$ verifying $((x_j')_j \neq x_j) \supset ((y_j')_j \neq y_j)$.

**Proof of the lemma.** Suppose that $z_j T_j y_j$ for some $z_j \in S(x_j)$. We have:

\[(x_j')_j \neq x_j \supset ((z_j')_j \neq z_j) \quad \text{for some } (z_j')_j \not \in X_j \times X_j.

(4.2) implies $(x_j')_j \neq x_j \supset ((z_j')_j \neq z_j) \not \in X_j \times X_j$ which contradicts $y_j \not V_j x_j$. The other part of the implication is established as follows:

Suppose that not $y_j \not V_j x_j$. Then: $(x_j')_j \neq x_j \supset ((y_j')_j \neq y_j)$ for some $(x_j')_j \neq x_j$, $(y_j')_j \neq y_j \not \in X_j \times X_j$ and thus $y_j \in S(x_j)$. As not $y_j T^4_j y_j$, this completes the proof.

**Proof of Theorem 4.1.**

(1) (4.1) and (4.2) $\Rightarrow$ $V_j$ interval order (by Lemma 4.4). Thus all we have to prove is that, $\forall j \in X_j, z_j \in X_j$:

\[
\begin{align*}
&x_j V_j y_j \quad \text{and} \quad z_j \not V_j w_j \quad \text{or} \quad \text{implies } x_j V_j w_j.
\end{align*}
\]

By Lemma 4.5, we have:

\[
\begin{align*}
x_j T^4_j y_j \quad \text{for all } y_j \in S(y_j) \quad \text{and} \quad z_j T^4_j w_j \quad \text{for all } w_j \in S(w_j).
\end{align*}
\]

Suppose now that not $(x_j V_j w_j)$ so that $w_j T_j x_j$ for some $w_j \in S(w_j)$. As $z_j \not T^4_j w_j$ and $z_j T^4_j y_j$ for all $y_j \in S(y_j)$, we have $z_j \not T^4_j y_j$ for all $y_j \in S(y_j)$, $T^4_j$ being a weak order. Thus $z_j \not V_j y_j$.

(2) (4.1), (4.2) and (4.3) $\Rightarrow$ $V_j$ semi-order. Given the first part of the theorem, all we have to prove is that $\forall x_j, y_j, z_j \in X_j$:

\[
\begin{align*}
&x_j V_j y_j \quad \text{and} \quad y_j \not V_j z_j \quad \text{implies } \exists \, w_j \in X_j \quad \text{or} \quad w_j \not V_j z_j.
\end{align*}
\]
First suppose that \( w_j T_j y_j \). As \( y_j V_j z_j \), \( y_j T_j^A z'_j \) for all \( z'_j \in S(z_j) \). Thus \( w_j T_j^A z'_j \) for all \( z'_j \in S(z_j) \) and \( w_j V_j z_j \).

Suppose now that \( y_j T_j^A w_j \). For any \( w_j \in S(w_j) \), we have \((w_j, (w_j^*)_{x,j}) \succ (\bar{w}_j, (w_j^*)_{x,j}) \) for some \((w_j^*)_{x,j} \in X_i \times X_i \) and by (4.3) \((y_j, (w_j^*)_{x,j}) \succ (w_j, (w_j^*)_{x,j}) \) which implies that \( \bar{w}_j \in S(y_j) \) for all \( w_j \in S(w_j) \). We have \( x_j V_j y_j \); therefore \( x_j T_j^A y_j' \) for all \( y_j' \in S(y_j) \). As \( S(w_j) \subset S(y_j) \), we have \( x_j V_j w_j \).

Theorem 4.1. is particularly useful because it generally allows (when there is a numerical representation of the relation \( V_j \)) to define for each \( x_j \) a "veto threshold".

The reader may be puzzled by the dissymmetry existing between (4.2) and (4.3), (4.3) being slightly weaker than (4.2). Hereafter, we study the consequences of replacing (4.3) by:

\[
\begin{align*}
\forall j & \in \Omega, \forall z_j \in X_j \text{ such that } \\
& z_j T_j X_j^* ((x_j^*)_{x,j}, z_j) \succ [y_j].
\end{align*}
\]

Let us first observe that, denoting \( U_j \) the transitive and strongly connected relation underlying \( V_j \), neither \( U_j \subset T_j \) nor \( T_j \subset U_j \) is a logical consequence of Theorem 4.1. When \( V_j \) is empty, it is obvious that \( U_j \subset T_j \) does not hold. In order to prove that we may not have \( T_j \subset U_j \) let us consider the following example of GNPS:

\[
\begin{align*}
X &= X_1 \times X_2 \\
X_1 &= \{x_1, y_1, z_1\} \\
X_2 &= \{x_2, y_2\}
\end{align*}
\]

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We have:
\[ x_1 \succ_1 y_1, y_1 \prec_1 z_1, x_1 \succ_1 z_1 \]
\[ x_2 \succ_2 y_2, \quad [1] \succ \emptyset, [2] \succ \emptyset \]
\[ [1], [2] \succ \emptyset \]

It can be verified that this GNPS is discordant, that each \( \succ \) is a weak-order and that \( \succ \) satisfies (4.2), (4.3) but not (4.4) because:
\( (y_1', x_2') \succ (x_1, y_2') \) and \( (z_1, x_2') \succ (x_1, y_2') \), though \( z_1 \mspace{1mu} T_{y_1} \mspace{1mu} y_1 \).

In this example, \( V_2 \) is empty and we have \( x_1 \mspace{1mu} V_1 \mspace{1mu} z_1 \). We have therefore
\[ x_1 \mspace{1mu} U_1 \mspace{1mu} z_1, \mspace{1mu} x_1 \mspace{1mu} U_1 \mspace{1mu} y_1 \text{ and } y_1 \mspace{1mu} U_1 \mspace{1mu} z_1 \]
but, at the same time
\[ x_1 \mspace{1mu} T_1 \mspace{1mu} z_1, \mspace{1mu} x_1 \mspace{1mu} T_1 \mspace{1mu} y_1, \mspace{1mu} y_1 \mspace{1mu} T_1 \mspace{1mu} z_1 \text{ and } z_1 \mspace{1mu} T_1 \mspace{1mu} y_1. \]

We complete our study of discordant GNPS by establishing:

**Theorem 4.2.** In a discordant GNPS verifying (4.1), (4.2) and (4.4), we have \( T_j \cup U_j \forall j \in \Omega \).

Remark before proof. \( T_j \) and \( U_j \) being strongly connected \( T_j \subseteq U_j \) implies \( I_{T_j} \subseteq U_{T_j} \) and \( U_{T_j} \subseteq I_{T_j} \).

Proof of Theorem 4.2. Suppose \( y_j \mspace{1mu} T_j \mspace{1mu} x_j \) and not \( (y_j \mspace{1mu} U_j \mspace{1mu} x_j) \). Not \( (y_j \mspace{1mu} U_j \mspace{1mu} x_j) \) implies:

- either \( z_j \mspace{1mu} V_j \mspace{1mu} y_j \) and not \( z_j \mspace{1mu} V_j \mspace{1mu} x_j \) for some \( z_j \in X_j \)
- or \( x_j \mspace{1mu} V_j \mspace{1mu} w_j \) and not \( y_j \mspace{1mu} V_j \mspace{1mu} w_j \) for some \( w_j \in X_j \).

Not \( (z_j \mspace{1mu} V_j \mspace{1mu} x_j) \Rightarrow (x_j, (x_i')_{i \neq j}) \succ (z_j, (z_i')_{i \neq j}) \) for some \( (x_i')_{i \neq j}, (z_i')_{i \neq j} \in X_{i \neq j} \).

(4.4) therefore implies:
\[ (y_j, (x_i')_{i \neq j}) \succ (z_j, (z_i')_{i \neq j}) \]

which contradicts \( z_j \mspace{1mu} V_j \mspace{1mu} y_j \).

Not \( (y_j, V_j \mspace{1mu} w_j) = (w_j, (w_i')_{i \neq j}) \succ (y_j, (y_i')_{i \neq j}) \) for some \( (w_i')_{i \neq j}, (y_i')_{i \neq j} \in X_{i \neq j} \).

and the application of (4.2) contradicts \( x_j \mspace{1mu} V_j \mspace{1mu} w_j \)
\[ \blacksquare \]
5. GNPS AND DECISION-AID

We mentioned in the introduction that the concept of noncompensation has been used in several multi-attribute decision-aid methods. Our purpose in this section is to outline how the theoretical considerations developed above can be helpful in order to design a new method using these concepts (the TACTIC method) and to implement it.

The idea of the TACTIC method is to “build” a global preference relation (see Bouyssou, 1984; and Roy and Bouyssou, 1986 on this notion of construction) given a set of several attributes. Technically, the global preference relation takes the form of a discordant GNPS verifying the conditions (4.1), (4.2) and (4.4), which associated NPS is a CPSp. Following Vansnick (1986), we will call such P.S. “noncompensatory preference structures with veto”. The method first seeks to determine, in agreement with the decision-maker, the semi-orders \( \succ_i \) and \( V_i \) for each attribute \( i \), with \( V_i \subsetneq \succ_i \). This can be done simply by assessing a measurable value function on each attribute determination two constant thresholds. The method then asks the decision-maker to compare several simple actions in order to obtain inter-criteria information. From this information, it determines simultaneously, following a number of reasonable principles, the ‘weights’ \( f_i \) and the coefficient \( p \). Once this information is obtained, the determination of \( \succ \) for the complete set of actions is performed easily. For more details on this topic, we refer to Vansnick (1986).

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