

# A Scenario Based Single Vehicle Routing Problem with Stochastic Demands: Flow Models

**Marcus Poggi.\***, Georges Spyrides, Thibaut Vidal

Departamento de Informática, PUC-Rio, Brazil

{ `poggi`, `gspyrides`, `vidalt` }@inf.puc-rio.br

\* and Accenture

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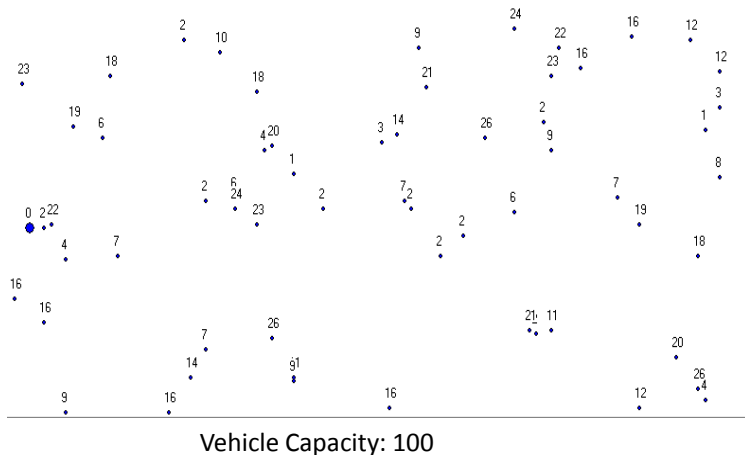
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  - VRP with Stochastic Demands
- 2 Motivation
- 3 Formulation
- 4 Experiments
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# Capacitated Vehicle Routing Problem (CVRP)

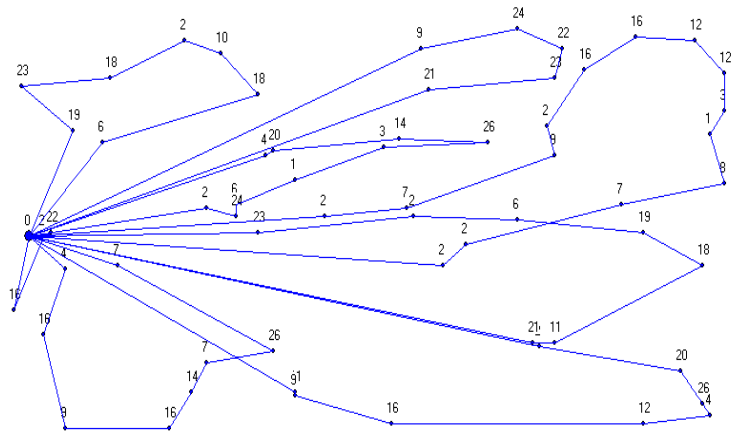
Given

- ☐ An undirected graph  $G=(V,E)$      $V= \{0,1,...,n\}$ 
  - Vertex 0 represents a depot and remaining vertices represent clients
  - Edge lengths are denoted by  $c(e)$
- ☐ Client demands are  $q(1),...,q(n)$
- ☐  $K$  vehicles with capacity  $C$
- ☐ Determine routes for each vehicle satisfying the following constraints:
  - (i) each route starts and ends at the depot,
  - (ii) each client is visited by exactly one vehicle
  - (iii) the total demand of clients visited in a route is at most  $C$
- ☐ The objective is to minimize the total route length.

# Example - Instance A-n62-k8



# Optimal Solution A-n62-k8



Optimal Solution Value: 1288

## Routing Problems with Stochastic Demands

- The common version assumes the demand at a client is only revealed at arrival: Pick-up problem
- The objective is to minimize the expected total distance traveled.
- Distances are deterministic (no probabilistic duration here)
- A **failure** along a route occurs when the client demand uncovered does not fit in the vehicle
- Different ways to manage **failures** (recourse) lead to different VRPSD versions
- Different ways to model the demand uncertainty lead to different approaches

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## Stochastic Demands: Some Literature

- Tillman, The Multiterminal delivery Problem with Stochastic Demands, TS, 1969
- Vehicle Routing with Stochastic Demands: Properties and Solution Frameworks Dror, Laporte, Trudeau, TS, 1989 (MDP)
- A vehicle routing problem with stochastic demand, Bertsimas, OR, 1992.

## Stochastic Demands: Managing Failures

- Recourse Action on client loading failure:
  - Preventive: Goes back to unload whenever there is a high probability of failure on the next client
  - Non-Preventive: occurs a failure on current client
    - Goes to the depot and returns to current client (popular!)
    - Goes to the depot and exchanges next, 2, 3, ... clients in an optimal way
    - Goes to the depot and reshuffles all remaining clients optimally
  - Chance constraints: guarantees a low probability of failure
  - Penalty for not serving a client
  - Optimal Re-stocking Policy: Given a fixed sequence of visits and demand distributions(i.i.) optimal re-stocking can be determined: for return to the depot and back to current client.

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## Stochastic Demands: More literature

- Louveaux, Salazar Gonzalez, 2018), arbitrary discrete distribution, preventive, small demands, L-Shaped
- Dinh, Fukasawa, Luedtke, 2017, Chance constraints
- Florio, Hartl, Minner, 2019, BCP (Probabilistic Duration), Optimal Re-stocking Policy (Dynamic Programming + BCP), non-preventive,
- Salavati, Gendreau, Jabali, Rei, 2019, Optimal Re-stocking Policy, Preventive and Non-preventive, L-Shaped

## Single Vehicle Routing with Stochastic Demands

- Zhang et al., T. Sci., 2014: Approx Dyn. Prog.
- Florio, Hartl, Minner, EJOR, 2018, MDP, Hamiltonian Policy

## Routing Problems with Stochastic Demands

- What applied VRPSD problems are out there?
- What experiment would supply evidence that the Stochastic approach should be chosen?
  - Take a (representative) history of working (similar) days demands, and simulate the proposed routing strategy
- How much would it help should there be long routes?
- How much uncertainty would have significant impact?
  - A first guess: Either some demands vary a lot, or there are positive correlations
  - Is the i.i.d. hypothesis adequate?
- There many transportation companies with steady pick-ups
- Many clients have significant uncertainty on the demand
- Mostly they execute the same routes every day

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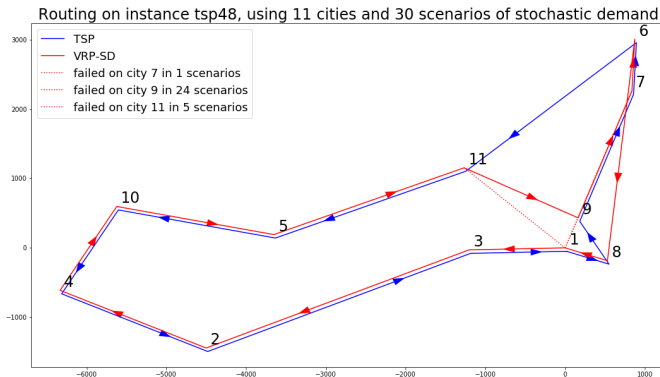
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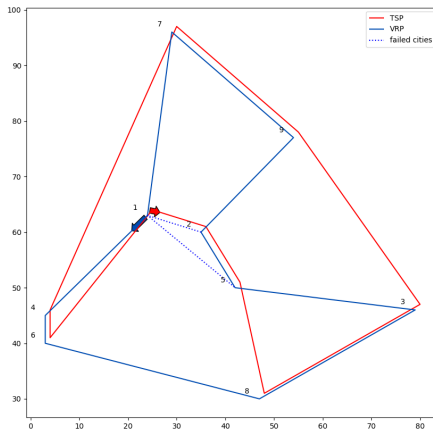
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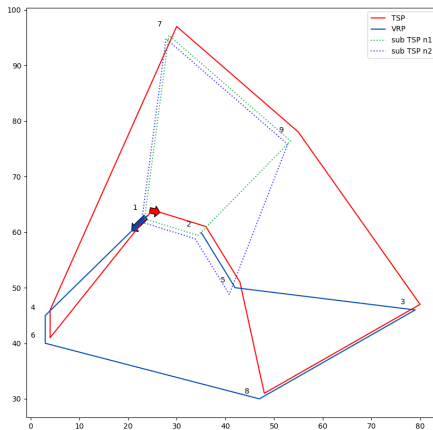
## Go to Depot and Back Recourse Policy



## Go to Depot and Back Recourse Policy



## Reshuffle Recourse Policy



## The VRPSDS Problem

### Scenario approach for Single Vehicle Routing

Given:

- A vehicle with capacity  $Q$ , travel costs  $c_{ij}$ , and demands  $d_{vs}$ ;
- Assume the sum of the demands fall in  $(Q, 2Q)$  with high probability
- In a given demand scenario a recourse must be taken on the failure client;
- Recourse function to manage *failure*:
  - Go to the depot unload and return to current client
  - Go to the depot and follow the optimal route for remaining clients

Determine:

- Sequence of client visits that minimizes the expected total traveling cost.

Optimal sequence is such that, for the scenarios considered, the expected overhead due to exceeding vehicle capacity is minimized

## The VRPSDS Problem

### Scenario approach OR Periodic VRP

- Periodic VRP: demands are known for the next  $S$  days
- Scenario approach: a sample of client demands from several "same days" are available

## Formulation Based on the Directed TSP Flow formulation

Let:

- $y_{ij}$ :  $(i, j)$  is used in the first level tour
- $x_{ijk}$ :  $(i, j)$  is traversed by the flow from vertex 1 to vertex  $k$

## Minimizing the Expected Total Distance Traveled

Single Vehicle TSP with Recourse Cost  $\alpha(y, x)$

$$\min \sum_{i=1}^n \sum_{j=2}^{n+1} c_{ij} y_{i,j} + \alpha(y, x)$$

TSP-Flow Constraints



## Minimizing the Expected Total Distance Traveled

Single Vehicle TSP with Recourse Cost  $\alpha(y, x)$ 

## TSP-Flow Constraints

$$\text{(out d)} \quad \sum_{i=1, i \neq j}^n y_{ij} = 1 \quad \forall j = 2, n+1$$

$$\text{(in d)} \quad \sum_{j=2, j \neq i}^{n+1} y_{ij} = 1 \quad \forall i = 1, n$$

$$\text{(flow out)} \quad \sum_{j=2}^n x_{1jk} = 1 \quad \forall k = 2, n+1$$

$$\text{(flow cons)} \quad \sum_{j=2}^{n+1} x_{ijk} - \sum_{j=1}^n x_{jik} = \begin{cases} -1 & i = k; \\ 0 & i \neq k \end{cases} \quad \forall i = 1, n \quad \forall k = 1, n$$

$$\text{(couple)} \quad \begin{aligned} x_{ijk} - y_{ij} &\leq 0 \\ y_{ij} &\in \{0, 1\} \\ x_{ijk} &\in \{0, 1\} \end{aligned} \quad \forall i = 1, n; \quad j = 2, n+1; \quad k = 2, n+1$$

Recourse: Return to the Depot and Back to the current client

$$\min \sum_{i=1}^n \sum_{j=2}^{n+1} c_{ij} y_{i,j} + \frac{1}{nscen} \sum_{s=1}^{nscen} \sum_{k=2}^n (c_{k(n+1)} + c_{1k}) \cdot pay_{ks}$$

TSP-Flow Constraints

Scenario Constraints

## Scenario Variables

Let:

- $o_i$ : contains the position of vertex  $i$  on the first level route
- $w_{is}$ : indicates client  $i$  is served after capacity failure in scenario  $s$ .
- $p_{ks}^+$ : the failure occurs after client  $k$  in the route
- $p_{ks}^-$ : if it is positive the failure occurs before client  $k$  in the route
- $s_{ks}$ : one if only  $p_{ks}^+$  can be positive, if it is zero, only  $p_{ks}^-$  can be positive
- $pay_{ks}$ : indicates fail on client  $k$  in scenario  $s$ : go to the depot and return
- $unpay_{ijs}$ : indicates the reshuffle recourse will not use edge  $(i, j)$  in  $s$

## Scenario Constraints

$$\begin{array}{ll}
\text{(serve)} & \sum_{i=1}^n \sum_{j=2, j \neq i}^n d_{js} x_{ijk} - Q \cdot w_{ks} \leq Q \quad \forall k = 2, \dots, n \\
\text{(serve comp)} & \sum_{i=1}^n \sum_{j=2, j \neq i}^n d_{js} x_{ijk} - Q \cdot w_{ks} \geq 0 \quad \forall k = 2, \dots, n \\
\text{(ord rout)} & \sum_{i=1}^n \sum_{j=2, j \neq i}^n x_{ijk} - o_k = 0 \quad \forall k = 2, \dots, n \\
\text{(prec ik)} & w_{is} - w_{ks} \leq 1 - x_{ijk} \quad \forall i > 1, i \neq k, k = 2, \dots, n \\
\text{(prec jk)} & w_{js} - w_{ks} \leq 1 - x_{ijk} \quad \forall j > 1, j \neq k, k = 2, \dots, n \\
\text{(2stage cl)} & \sum_{k=2}^n d_{ks} w_{ks} \leq Q \quad \forall s \\
\text{(notserv)} & \sum_{k=2}^n w_{ks} - n n s_s = 0 \quad \forall s \\
\text{(last serv)} & o_k + n n s_s + p_{ks}^+ - p_{ks}^- = n \quad \forall k = 2, \dots, n \quad \forall s \\
\text{(pos sign)} & p_{ks}^+ - n \cdot s_{ks} \leq 0 \quad \forall k = 2, \dots, n \quad \forall s \\
\text{(neg sign)} & p_{ks}^- + n \cdot s_{ks} \leq n \quad \forall k = 2, \dots, n \quad \forall s \\
\text{(set pay)} & p a y_{ks} + p_{ks}^+ + p_{ks}^- \geq 1 \quad \forall k = 2, \dots, n \quad \forall s
\end{array}$$

## Reshuffle Recourse

Return to the depot and perform optimal route on clients not served

$$\min \sum_{i=1}^n \sum_{j=2}^{n+1} c_{ij} y_{i,j} + \frac{1}{nscen} \sum_{s=1}^{nscen} \left( \sum_{k=2}^n c_{k(n+1)} \cdot pay_{ks} - \sum_{i=1}^n \sum_{j=2}^{n+1} c_{ij} \cdot unpaid_{ijs} + \sum_{i=1}^n \sum_{j=2}^{n+1} c_{ij} \cdot g_{ijs} \right)$$

## TSP-Flow Constraints

## Scenario Constraints

## TSP Scenario

## TSP Scenario

$$\sum_{i=1}^n g_{ijs} - w_{js} = 0$$

$$\forall j = 2, \dots, n+1$$

$$\sum_{j=2}^{n+1} g_{ijs} - w_{is} = 0$$

$$\forall i = 1, \dots, n$$

$$\sum_{j=2}^n f_{1jks} - w_{ks} = 0$$

$$\forall k = 2, \dots, n+1$$

$$\sum_{j=2}^{n+1} f_{ijks} - \sum_{j=1}^n f_{jiks} = \{-w_{ks} \mid i = k; \ 0 \mid i \neq k\}$$

$$\forall i = 1, \dots, n \quad \forall k = 1, \dots, n$$

$$unpay_{ijs} \leq w_{is}$$

$$unpay_{ijs} \leq w_{js}$$

$$unpay_{ijs} \geq w_{is} + w_{js} - 1$$

$$unpay_{ijs} \in \{0, 1\}$$

$$g_{ijs} \in \{0, 1\}$$

$$f_{ijks} \in \{0, 1\}$$

## Combinatorial Benders Cuts

Let  $y^+$  be a solution for (TSP constraints)

Let  $zr_\alpha(y^+)$  be the expected recourse cost associated with  $y^+$

Then:

$$\alpha \geq zr_\alpha(y^+).(1 - n + \sum_{(i,j)|y_{ij}^+=1} y_{ij})$$

is the associated Benders Combinatorial (or Optimality) cut.

## Prefix Combinatorial Benders Cuts

Let  $y^p$  be the incidence vector of edges corresponding to a **prefix** of a route

Let  $p$  be the number of edges in the path starting at the depot

Let  $zr_\alpha^s(y^p)$  be the recourse cost associated with  $y^p$  in scenario  $s$  (Depot and back or Reshuffle)

Let also  $\alpha_s$  be the recourse cost associated to scenario  $s$

If in scenario  $s$  we have  $\sum_{(i,j)|y_{ij}^p=1} d_{js} > Q$ ; Then:

$$\alpha_s \geq zr_\alpha^s(y^p) \cdot (1 - p + \sum_{(i,j)|y_{ij}^p=1} y_{ij})$$

is the associated Benders Combinatorial (or Optimality) cut for scenario  $s$ .



## Naive 1: Benders Decomposition - Two Stage

- Initialization: set  $i = 1$ , a lower bound  $\underline{z} = -\infty$ , and an upper bound  $\bar{z} = +\infty$ 
  - S1 Solve current first stage determining trial solution  $y^i$  and let  $\underline{z} = z(y^i)$
  - S2 Solve the second stage determining  $zr_\alpha(y^i)$
  - S3 Set  $\bar{z} = c.y^i + zr_\alpha(y^i)$
  - S4 Add the Benders Combinatorial (or Optimality) Cut to the first stage.
  - S5 IF  $\bar{z} - \underline{z} < \epsilon$ . STOP.  
Otherwise set  $i = i + 1$ , goto S1

## Algorithms tested

- Naive Benders with Combinatorial Cuts
- Proposed model as is.

## Instances

- Based on TSPLIB instances. Euclidean instances are hard.
- Demands:
  - Mean  $\mu_i$  is generated as UNIF[50, 80] for each client
  - Client  $i$  in scenario  $s$  has demand  $q_{is}$  from UNIF[ $0.7\mu_i, 1.3\mu_i$ ]

We also experiment with more uneven demands.

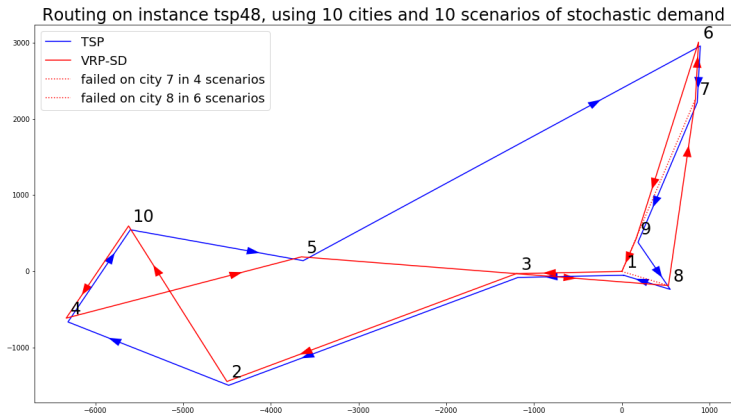
## Test-Bed

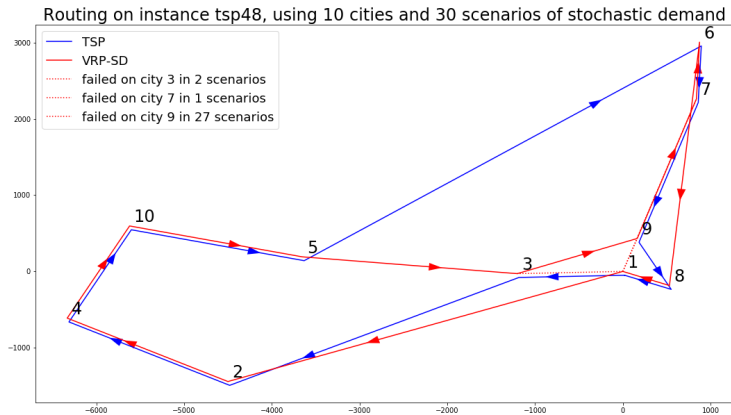
- All computational experiments were executed on the cloud, on VMs with the following characteristics:
  - Quad core intel Xeon processor @2.6 Ghz
  - 26 Gb of RAM
  - Windows Server 2012 - 64bit
- Code developed using Julia 1.1 + JuMP modeling language.
- Mathematical programming solver: Gurobi

## Results for Recourse Policy 1: Go to the depot and back

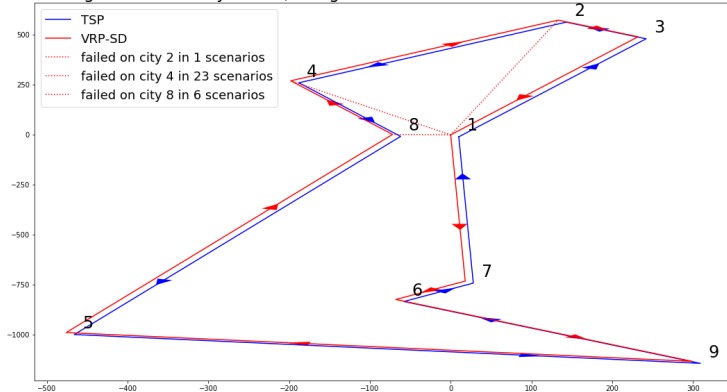
- TSP with combinatorial Benders can solve until 12 in a few minutes
- The model also solves for 12, takes longer
- Integrality gap between 5% and 15%
- Sensitive to instances (very easy and very hard)
- Large recourse cost implies harder instances
- Largest instance solved has 15 clients

We present next some optimal solutions





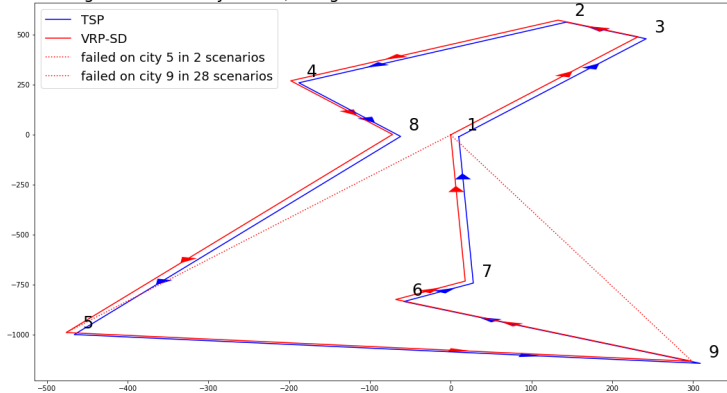
Routing on instance ulysses16, using 9 cities and 30 scenarios of stochastic demand



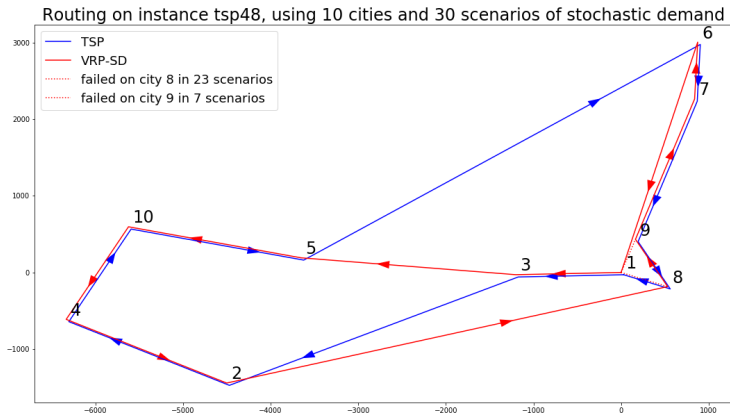


	2	3	4	5	6	7	8	9
0	92.0	46.0	79.0	47.0	45.0	63.0	66.0	49.0
1	66.0	80.0	60.0	46.0	68.0	49.0	59.0	62.0
2	68.0	51.0	70.0	69.0	75.0	56.0	49.0	77.0
3	79.0	47.0	64.0	50.0	58.0	70.0	47.0	50.0
4	90.0	45.0	64.0	45.0	44.0	48.0	58.0	58.0
5	70.0	65.0	83.0	64.0	61.0	59.0	49.0	64.0
6	55.0	55.0	68.0	47.0	60.0	62.0	84.0	63.0
7	80.0	69.0	60.0	79.0	44.0	77.0	52.0	74.0
8	81.0	70.0	59.0	48.0	65.0	63.0	62.0	54.0
9	92.0	61.0	92.0	66.0	42.0	54.0	82.0	83.0
10	67.0	61.0	94.0	78.0	66.0	83.0	66.0	54.0
11	83.0	61.0	92.0	63.0	75.0	48.0	47.0	56.0
12	84.0	62.0	69.0	51.0	74.0	56.0	60.0	79.0
13	80.0	69.0	69.0	59.0	64.0	71.0	68.0	70.0
14	91.0	75.0	65.0	49.0	56.0	76.0	52.0	59.0
15	81.0	71.0	74.0	80.0	45.0	76.0	47.0	67.0
16	99.0	63.0	86.0	71.0	60.0	66.0	65.0	87.0
17	63.0	56.0	101.0	45.0	42.0	48.0	65.0	77.0
18	85.0	59.0	64.0	63.0	73.0	48.0	78.0	83.0
19	66.0	57.0	99.0	64.0	61.0	75.0	66.0	66.0
20	99.0	48.0	71.0	47.0	42.0	52.0	74.0	70.0
21	78.0	74.0	59.0	66.0	69.0	65.0	81.0	63.0
22	69.0	53.0	57.0	59.0	72.0	53.0	72.0	71.0
23	58.0	72.0	77.0	79.0	42.0	68.0	62.0	52.0
24	57.0	45.0	86.0	78.0	62.0	75.0	54.0	47.0
25	85.0	57.0	92.0	52.0	60.0	64.0	69.0	84.0
26	52.0	48.0	61.0	57.0	68.0	80.0	80.0	77.0
27	78.0	67.0	82.0	83.0	65.0	76.0	75.0	55.0
28	54.0	66.0	91.0	95.0	70.0	58.0	60.0	60.0
29	71.0	48.0	90.0	68.0	60.0	78.0	64.0	56.0

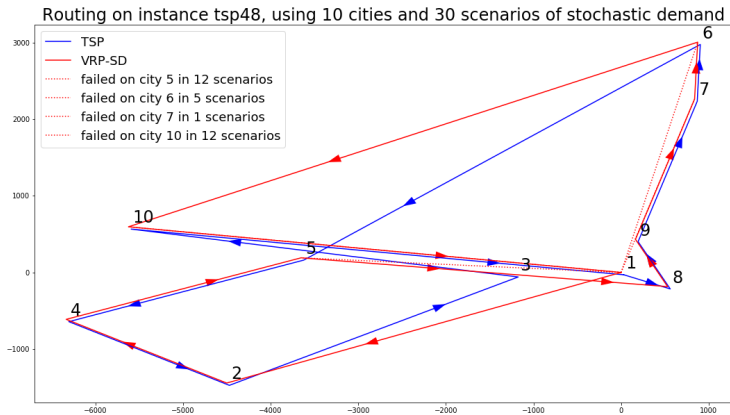
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1	34.0	36.0	28.0	77.0	48.0	48.0	28.0	70.0
2	39.0	31.0	30.0	71.0	29.0	44.0	16.0	65.0
3	24.0	27.0	26.0	66.0	37.0	52.0	21.0	88.0
4	33.0	33.0	23.0	63.0	45.0	47.0	20.0	97.0
5	24.0	30.0	18.0	65.0	36.0	31.0	21.0	100.0
6	41.0	40.0	30.0	88.0	35.0	38.0	27.0	75.0
7	26.0	39.0	30.0	77.0	46.0	50.0	24.0	99.0
8	30.0	30.0	22.0	57.0	44.0	32.0	24.0	59.0
9	26.0	23.0	24.0	74.0	40.0	43.0	25.0	80.0
10	30.0	33.0	28.0	66.0	48.0	47.0	19.0	67.0
11	36.0	31.0	29.0	76.0	31.0	33.0	27.0	87.0
12	30.0	30.0	30.0	52.0	32.0	50.0	25.0	100.0
13	40.0	27.0	30.0	89.0	29.0	43.0	25.0	65.0
14	40.0	23.0	30.0	65.0	42.0	28.0	17.0	70.0
15	38.0	38.0	29.0	55.0	39.0	34.0	28.0	66.0
16	28.0	32.0	22.0	57.0	42.0	45.0	25.0	84.0
17	32.0	23.0	27.0	58.0	42.0	32.0	19.0	75.0
18	27.0	26.0	24.0	60.0	42.0	39.0	18.0	99.0
19	31.0	22.0	31.0	58.0	37.0	36.0	27.0	88.0
20	25.0	34.0	24.0	51.0	35.0	37.0	24.0	89.0
21	28.0	25.0	28.0	81.0	44.0	35.0	24.0	60.0
22	39.0	28.0	21.0	63.0	37.0	30.0	26.0	64.0
23	41.0	30.0	20.0	76.0	44.0	51.0	24.0	97.0
24	26.0	33.0	27.0	58.0	44.0	46.0	22.0	65.0
25	33.0	31.0	22.0	65.0	28.0	43.0	24.0	66.0
26	30.0	28.0	18.0	70.0	26.0	40.0	25.0	72.0
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28	34.0	26.0	29.0	70.0	30.0	52.0	35.0	55.0
29	27.0	27.0	19.0	76.0	35.0	42.0	19.0	58.0

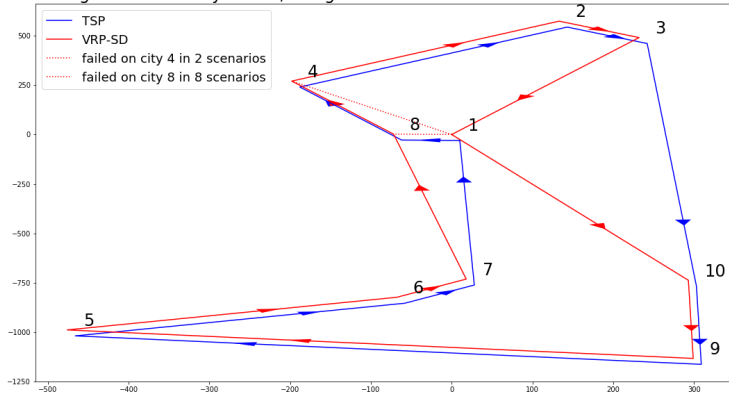


	2	3	4	5	6	7	8	9	10
0	89.0	67.0	49.0	72.0	76.0	79.0	66.0	48.0	64.0
1	61.0	56.0	46.0	91.0	74.0	61.0	71.0	50.0	50.0
2	58.0	78.0	54.0	76.0	82.0	61.0	73.0	54.0	46.0
3	81.0	56.0	56.0	89.0	47.0	73.0	70.0	68.0	63.0
4	71.0	54.0	52.0	61.0	82.0	61.0	50.0	61.0	53.0
5	54.0	50.0	56.0	82.0	61.0	72.0	66.0	64.0	60.0
6	52.0	55.0	42.0	62.0	51.0	48.0	71.0	69.0	52.0
7	53.0	45.0	62.0	73.0	80.0	59.0	66.0	53.0	52.0
8	79.0	76.0	46.0	78.0	73.0	74.0	59.0	63.0	48.0
9	62.0	70.0	43.0	54.0	66.0	56.0	55.0	62.0	64.0
10	68.0	59.0	56.0	77.0	74.0	67.0	86.0	61.0	36.0
11	70.0	62.0	36.0	82.0	84.0	61.0	54.0	66.0	44.0
12	67.0	58.0	57.0	85.0	53.0	65.0	53.0	70.0	62.0
13	56.0	58.0	65.0	64.0	83.0	35.0	57.0	52.0	61.0
14	80.0	43.0	40.0	50.0	73.0	72.0	58.0	66.0	53.0
15	59.0	50.0	55.0	58.0	64.0	66.0	62.0	52.0	45.0
16	65.0	71.0	50.0	90.0	63.0	66.0	73.0	54.0	45.0
17	90.0	64.0	46.0	61.0	77.0	75.0	62.0	46.0	48.0
18	67.0	52.0	43.0	81.0	61.0	65.0	59.0	44.0	59.0
19	78.0	68.0	53.0	68.0	74.0	76.0	51.0	51.0	37.0
20	52.0	75.0	49.0	86.0	74.0	63.0	57.0	75.0	52.0
21	74.0	78.0	63.0	64.0	46.0	56.0	62.0	68.0	37.0
22	54.0	71.0	66.0	68.0	83.0	77.0	50.0	52.0	63.0
23	69.0	62.0	43.0	55.0	56.0	68.0	65.0	48.0	38.0
24	59.0	58.0	53.0	67.0	55.0	44.0	86.0	42.0	60.0
25	59.0	58.0	40.0	55.0	77.0	59.0	57.0	55.0	45.0
26	74.0	45.0	63.0	52.0	55.0	74.0	81.0	75.0	64.0
27	83.0	60.0	52.0	51.0	63.0	49.0	69.0	45.0	41.0
28	76.0	70.0	39.0	56.0	56.0	78.0	55.0	45.0	35.0
29	63.0	78.0	52.0	53.0	49.0	58.0	71.0	70.0	53.0

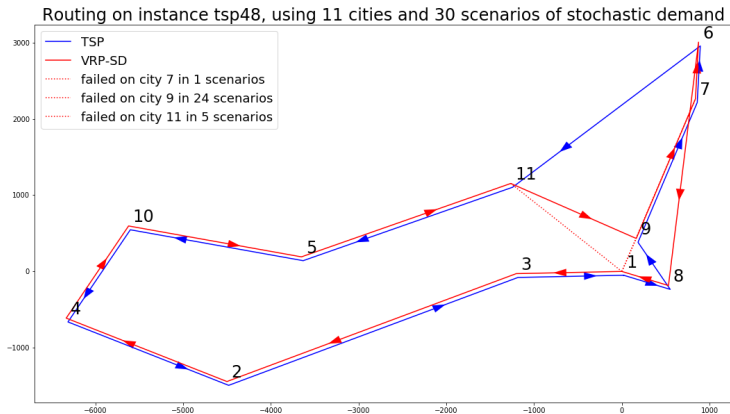


	2	3	4	5	6	7	8	9	10
0	39.0	26.0	55.0	31.0	39.0	35.0	17.0	25.0	71.0
1	39.0	20.0	92.0	32.0	29.0	34.0	20.0	27.0	73.0
2	32.0	19.0	81.0	33.0	39.0	34.0	20.0	31.0	59.0
3	48.0	22.0	78.0	39.0	31.0	24.0	23.0	28.0	52.0
4	51.0	30.0	88.0	29.0	29.0	27.0	24.0	17.0	70.0
5	39.0	24.0	87.0	27.0	32.0	36.0	19.0	23.0	67.0
6	55.0	20.0	88.0	35.0	30.0	25.0	23.0	20.0	60.0
7	36.0	20.0	60.0	30.0	29.0	27.0	21.0	18.0	57.0
8	38.0	28.0	64.0	22.0	35.0	35.0	14.0	30.0	75.0
9	50.0	25.0	83.0	22.0	30.0	32.0	17.0	18.0	60.0
10	57.0	32.0	54.0	35.0	35.0	23.0	16.0	26.0	76.0
11	38.0	23.0	81.0	21.0	30.0	36.0	14.0	19.0	67.0
12	43.0	20.0	63.0	31.0	24.0	29.0	16.0	18.0	63.0
13	53.0	20.0	80.0	33.0	22.0	39.0	18.0	19.0	75.0
14	49.0	18.0	57.0	36.0	25.0	34.0	19.0	23.0	44.0
15	53.0	18.0	77.0	21.0	40.0	26.0	25.0	27.0	53.0
16	35.0	26.0	85.0	26.0	33.0	32.0	26.0	17.0	52.0
17	50.0	32.0	59.0	28.0	26.0	33.0	17.0	19.0	55.0
18	53.0	29.0	67.0	24.0	39.0	35.0	21.0	26.0	65.0
19	37.0	19.0	89.0	22.0	25.0	41.0	20.0	29.0	61.0
20	55.0	24.0	57.0	27.0	38.0	23.0	23.0	17.0	67.0
21	37.0	31.0	68.0	37.0	26.0	33.0	19.0	21.0	72.0
22	34.0	28.0	62.0	29.0	22.0	38.0	15.0	28.0	74.0
23	51.0	18.0	64.0	35.0	37.0	30.0	17.0	29.0	61.0
24	32.0	28.0	59.0	22.0	37.0	33.0	16.0	18.0	74.0
25	45.0	31.0	80.0	33.0	23.0	35.0	18.0	19.0	61.0
26	49.0	20.0	81.0	33.0	23.0	34.0	15.0	19.0	73.0
27	47.0	18.0	71.0	22.0	30.0	35.0	17.0	22.0	48.0
28	42.0	19.0	90.0	22.0	39.0	23.0	20.0	26.0	74.0
29	48.0	19.0	56.0	38.0	24.0	27.0	22.0	26.0	54.0

Routing on instance ulysses16, using 10 cities and 10 scenarios of stochastic demand







## Conclusions

- A first look at the problems shows difficulty starts for routes over 13 clients
- This means it is close to enumeration (matches the harder CVRP instances for BCP)
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### Some Potentially Non-Investigated Questions:

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## Future work

- Study lifting of classical Benders cuts (generating route is known)
- Collect and try real Stochastic Demands for VRP's

# Thank you!