# Cooperative Games Lecture 7: The Shapley Value

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# The Shapley value

Lloyd S. Shapley. A Value for n-person Games. In Contributions to the Theory of Games, volume II (Annals of Mathematical Studies), 1953.



**Definition** (marginal contribution)

The **marginal contribution** of agent *i* for a coalition  $\mathcal{C} \subseteq N \setminus \{i\}$  is  $mc_i(\mathcal{C}) = v(\mathcal{C} \cup \{i\}) - v(\mathcal{C})$ .

 $\langle mc_1(\emptyset), mc_2(\{1\}), mc_3(\{1,2\}) \rangle$  is an efficient payoff distribution for any game ( $\{1,2,3\}, v$ ). This payoff distribution may model a dynamic process in which 1 starts a coalition, is joined by 2, and finally 3 joins the coalition  $\{1,2\}$ , and where the incoming agent gets its marginal contribution.

An agent's payoff depends on which agents are already in the coalition. This payoff may not be **fair**. To increase fairness,one could take the average marginal contribution over all possible joining orders.

Let  $\sigma$  represent a joining order of the grand coalition *N*, i.e.,  $\sigma$  is a permutation of  $\langle 1, ..., n \rangle$ .

We write  $mc(\sigma) \in \mathbb{R}^n$  the payoff vector where agent *i* obtains  $mc_i(\{\sigma(j) | j < i\})$ . The vector *mc* is called a **marginal vector**.

Let (N, v) be a TU game. Let  $\Pi(N)$  denote the set of all permutations of the sequence  $\langle 1, \ldots, n \rangle$ .

$$Sh(N,v) = rac{\displaystyle\sum_{\sigma \in \Pi(N)} mc(\sigma)}{n!}$$

the Shapley value is a fair payoff distribution based on marginal contributions of agents averaged over joining orders of the coalition.

# An example

$\begin{split} N = &\{1,2,3\}, \ v(\{1\}) = 0, \ v(\{2\}) = 0, \ v(\{3\}) = 0, \\ &v(\{1,2\}) = 90, \ v(\{1,3\}) = 80, \ v(\{2,3\}) = 70, \\ &v(\{1,2,3\}) = 120. \end{split}$								
	1	2	3	Le	Let $y = (50, 40, 30)$			
$1 \leftarrow 2 \leftarrow 3$	0	90	30		C	$e(\mathcal{C}, x)$	$e(\mathcal{C},y)$	
$1 \leftarrow 3 \leftarrow 2$	0	40	80		{1}	-45	0	
$2 \leftarrow 1 \leftarrow 3$	90	0	30		{2}	-40	0	
$2 \leftarrow 3 \leftarrow 1$	50	0	70		{3}	-35	0	
$3 \leftarrow 1 \leftarrow 2$	80	40	0		{1,2}	5	0	
$3 \leftarrow 2 \leftarrow 1$	50	70	0		{1,3}	0	0	
total	270	240	210	•	{2,3}	-5	0	
Shapley value	45	40	35	{	1,2,3}	120	0	

This example shows that the Shapley value may not be in the core, and may not be the nucleolus.

- There are  $|\mathcal{C}|!$  permutations in which all members of  $\mathcal{C}$  precede *i*.
- There are  $|N \setminus (\mathcal{C} \cup \{i\})|!$  permutations in which the remaining members succede *i*, i.e.  $(|N| |\mathcal{C}| 1)!$ .

The Shapley value  $Sh_i(N, v)$  of the TU game (N, v) for player *i* can also be written

$$Sh_i(N,v) = \sum_{\mathcal{C} \subseteq N \setminus \{i\}} \frac{|\mathcal{C}|!(|N| - |\mathcal{C}| - 1)!}{|N|!} \left( v(\mathcal{C} \cup \{i\}) - v(\mathcal{C}) \right).$$

Using definition, the sum is over  $2^{|N|-1}$  instead of |N|!.

## **Definition** (value function)

Let  $\mathcal{G}_N$  the set of all TU games (N, v). A value function  $\phi$  is a function that assigns to each TU game (N, v) an efficient allocation, i.e.  $\phi : \mathcal{G}_N \to \mathbb{R}^{|N|}$  such that  $\phi(N, v)(N) = v(N)$ .

- The Shapley value is a value function.
- None of the concepts presented thus far were a value function (the nucleolus is guaranteed to be non-empty only for games with a non-empty set of imputations).

Let (N, v) and (N, u) be TU games and  $\phi$  be a value function.

- Symmetry or substitution (SYM): If  $\forall (i,j) \in N$ ,  $\forall \mathcal{C} \subseteq N \setminus \{i, j\}, v(\mathcal{C} \cup \{i\}) = v(\mathcal{C} \cup \{j\}) \text{ then } \phi_i(N, v) = \phi_i(N, v)$
- **Dummy (DUM):** If  $\forall \mathcal{C} \subset N \setminus \{i\} \ v(\mathcal{C}) + v(\{i\}) = v(\mathcal{C} \cup \{i\})$ , then  $\phi_i(N, v) = v(\{i\}).$
- Additivity (ADD): Let (N, u+v) be a TU game defined by  $\forall \mathcal{C} \subseteq N$ , (u+v)(N) = u(N) + v(N).  $\Phi(u+v) = \Phi(u) + \Phi(v).$

#### Theorem

The Shapley value is the unique value function  $\phi$  that satisfies (SYM), (DUM) and (ADD).

Let *N* be a set of agents and  $T \subseteq N \setminus \emptyset$ . The **unanimity game**  $(N, v_T)$  is defined as follows:  $\forall \mathcal{C} \subseteq N, v_T(\mathcal{C}) = \begin{cases} 1, \text{ if } T \subseteq \mathcal{C}, \\ 0 \text{ otherwise.} \end{cases}$ 

We note that

- if  $i \in N \setminus T$ , *i* is a null player.
- if  $(i,j) \in T^2$ , *i* and *j* are substitutes.

#### Lemma

The set  $\{v_T \mid T \subseteq N \setminus \emptyset\}$  is a linear basis of  $\mathcal{G}_N$ .

This means that a TU game (N, v) can be represented by a unique set of values  $(\alpha_T)_{T \subseteq N \setminus \emptyset}$  such that

$$\forall \mathfrak{C} \subseteq N, v(\mathfrak{C}) = \left(\sum_{T \subseteq N \setminus \emptyset} \alpha_T v_T\right)(\mathfrak{C}).$$

There are  $2^n - 1$  unanimity games and the dimension of  $\mathcal{G}_N$  is also  $2^n - 1$ .

We only need to prove that the unanimity games are linearly independent.

Towards a contradiction, let us assume that  $\sum_{T \subseteq N \setminus \emptyset} \alpha_T v_T = 0$ where  $(\alpha_T)_{T \subseteq N \setminus \emptyset} \neq 0_{\mathbb{R}^{2^n-1}}$ . Let  $T_0$  be a minimal set in  $\{T \subseteq N \mid \alpha_T \neq 0\}$ . Then,  $\left(\sum_{T \subseteq N \setminus \emptyset} \alpha_T v_T\right)(T_0) = \alpha_{T_0} \neq 0$ , which is a contradic-

tion.  $(\sum_{T \subseteq N \setminus \emptyset} \alpha_T \sigma_T)(T_0) = \alpha_{T_0} \neq 0$ , which is a contract

Let  $\phi$  a feasible solution on  $\mathcal{G}_N$  that is non-empty and satisfies the axioms SYM, DUM and ADD. Let us prove that  $\phi$  is a value function.Let  $(N, v) \in \mathcal{G}_N$ .

- if  $v = 0_{G_M}$ , all players are dummy. Since the solution is non-empty,  $0^{\mathbb{R}^{|N|}}$  is the unique member of  $\phi(N, v)$ .
- otherwise,  $(N, -v) \in \mathcal{G}_N$ . Let  $x \in \phi(N, v)$  and  $y \in \phi(N, -v)$ . By ADD,  $x + y \in \phi(v - v)$ , and then, x = -y is unique. Moreover,  $x(N) \leq v(N)$  as  $\phi$  is a feasible solution. Also  $y(N) \leq -v(N)$ . Since x = -y, we have  $v(N) \leq x(N) \leq v(N)$ , i.e. x is efficient.

Hence,  $\phi$  is a value function.

Let  $T \subseteq N \setminus \emptyset$  and  $\alpha \in \mathbb{R}$ . Let us prove that  $\phi(N, \alpha \cdot v_T)$  is uniquely defined.

- Let  $i \notin T$ . We have trivially  $T \subseteq C$  iff  $T \subseteq C \cup \{i\}$ . Then  $\forall C \subseteq N \setminus \{i\}, \ \alpha v_T(C) = \alpha v_T(C \cup \{i\})$ .Hence, all agent  $i \notin T$ are dummies. By DUM,  $\forall i \notin T, \ \phi_i(N, \alpha \cdot v_T) = 0$ .
- Let  $(i,j) \in T^2$ . Then for all  $\mathbb{C} \subseteq N \setminus \{i,j\}$ ,  $v(\mathbb{C} \cup \{i\}) = v(\mathbb{C} \cup \{j\})$ . By SYM,  $\phi_i(N, \alpha \cdot v_T) = \phi_j(N, \alpha \cdot v_T)$ .
- Since  $\phi$  is a value function, it is efficient. Then,  $\sum_{i \in N} \phi_i(N, \alpha \cdot v_T) = \alpha v_T(N) = \alpha.$ Hence,  $\forall i \in T$ ,  $\phi_i(N, \alpha \cdot v_T) = \frac{\alpha}{|T|}$ .

This proves that  $\phi(N, \alpha \cdot v_T)$  is uniquely defined. Since any TU game (N, v) can be written as  $\sum_{T \subseteq N \setminus \emptyset} \alpha_T v_T$  and because of ADD, there is a unique value function that satisfies the three axioms.

We need to show that the Shapley value satisfies the three axioms. Let (N, v) a TU game.

$$Sh(N,v) = rac{\displaystyle \sum_{\sigma \in \Pi(N)} mc(\sigma)}{n!}$$

• Let us assume that  $\forall C \subseteq N \setminus \{i, j\}$ , we have  $v(C \cup \{i\}) = v(C \cup \{j\})$ . Then  $\forall C \subseteq N \setminus \{i, j\}$ , we have

• 
$$mc_i(\mathcal{C}) = mc_j(\mathcal{C})$$

•  $v(\mathcal{C} \cup \{i,j\}) - v(\mathcal{C} \cup \{i\}) = v(\mathcal{C} \cup \{i,j\}) - v(\mathcal{C} \cup \{j\})$ , hence, we have  $mc_j(\mathcal{C} \cup \{j\}) = mc_i(\mathcal{C} \cup \{i\})$ .

$$\Rightarrow$$
  $Sh_i(N,v) = Sh_i(N,v)$ , Sh satisfies SYM.

- Let us assume there is an agent *i* such that for all  $C \subseteq N \setminus \{i\}$  we have  $v(C) = v(C \cup \{i\})$ . Then, each marginal contribution of player *i* is zero, and it follows that  $Sh_i(N, v) = 0$ . *Sh* satisfies DUM.
- *Sh* is clearly additive.

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## Discussion about the axioms

- SYM: it is desirable that two subsitute agents obtain the same value ✓
- DUM: it is desirable that an agent that does not bring anything in the cooperation does not get any value. ✓
- What does the addition of two games mean?
  - + if the payoff is interpreted as an expected payoff, ADD is a desirable property.
  - + for cost-sharing games, the interpretation is intuitive: the cost for a joint service is the sum of the costs of the separate services.
  - there is no interaction between the two games.
  - the structure of the game (N, v + w) may induce a behavior that has may be unrelated to the behavior induced by either games (N, v) or (N, w).
- The axioms are independent. If one of the axiom is dropped, it is possible to find a different value function satisfying the remaining two axioms.

Let (N, v) and (N, v) be two TU games.

• Marginal contribution: A value function  $\phi$  satisfies marginal contribution axiom iff for all  $i \in N$ , if for all  $\mathcal{C} \subset N \setminus \{i\}$   $v(\mathcal{C} \cup \{i\}) - v(C) = u(\mathcal{C} \cup \{i\}) - u(C)$ , then  $\phi(u) = \phi(v)$ .

The value of a player depends only on its marginal contribution.

# Theorem (H.P. Young)

The Shapley value is the unique value function that satisfies symmetry and marginal contribution axioms.

We refer by  $v \setminus i$  the TU game  $(N \setminus \{i\}, v_{\setminus i})$  where  $v_{\setminus i}$  is the restriction of *v* to  $N \setminus \{i\}$ .

• **Balanced contribution:** A value function  $\phi$  satisfies balanced contribution iff for all  $(i, j) \in N^2$  $\Phi_i(v) - \Phi_i(v \setminus j) = \Phi_j(v) - \Phi_j(v \setminus i).$ 

For any two agents, the amount that each agent would win or lose if the other "leaves the game" should be the same.

## Theorem (R Myerson)

The Shapley value is the unique value function that satisfies the balanced contribution axiom.

# Some properties

### Theorem

For superadditive games, the Shapley value is an imputation.

#### Lemma

For convex game, the Shapley value is in the core.



- Let (N, v) be a superadditive TU game. By superadditivity,  $\forall i \in N$ ,  $\forall C \subseteq N \setminus \{i\}$  $v(C \cup \{i\}) - v(C) > v(\{i\})$ . Hence, for each marginal vector, an agent *i* gets at least  $v(\{i\})$ . The same is true for the Shapley value as it is the average over all marginal vectors.
- Let (*N*, *v*) be a convex game.
  We know that all marginal vectors are in the core (to show that convex games have non-empty core, we used one marginal vector and showed it was in the core).
  The core is a convex set.
  The average of a finite set of points in a convex set is

also in the set.

Finally, the Shapley value is in the core.

# Summary

#### pros

- The Shapley value is a value function, i.e., it always exists and is unique.
- When the valuation function is **superadditive**, the Shapley value is **individually rational**, i.e., it is an imputation.
- When the valuation function is **convex**, the Shapley value is also group rational, hence, it is in the **core**.
- The Shapley value is the unique value function satisfying some axioms.

#### cons

• The nature of the Shapley value is combinatorial.



• Voting games and power indices.